

# Predicate Logic

## (SE 212 Tutorial 5)

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# Today's plan

- do some semantics questions from homework 4
- do some ND questions from homework 5

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- do some ND questions from homework 5



## h04q05 (cont'd)

Domain:

$$M = \{m1, m2\}$$

$$N = \{n1, n2\}$$

Mapping:

| Syntax |  | Meaning |
|--------|--|---------|
|--------|--|---------|

-----

|      |  |             |
|------|--|-------------|
| g(.) |  | G(n1) := m1 |
|      |  | G(n2) := m2 |

-----

|         |  |                |
|---------|--|----------------|
| p(., .) |  | P(m1, m1) := F |
|         |  | P(m1, m2) := T |
|         |  | P(m2, m1) := T |
|         |  | P(m2, m2) := F |

Premise:

$$\begin{aligned}
 & [\text{forall } y : M . \text{exists } x : N . p(g(x), y)] \\
 = & [\text{exists } x : N . p(g(x), \sim m1)] \text{ AND} \\
 & [\text{exists } x : N . p(g(x), \sim m2)] \\
 = & (P(G(n1), m1) \text{ OR } P(G(n2), m1)) \text{ AND} \\
 & (P(G(n1), m2) \text{ OR } P(G(n2), m2)) \\
 = & (P(m1, m1) \text{ OR } P(m2, m1)) \text{ AND} \\
 & (P(m1, m2) \text{ OR } P(m2, m2)) \\
 = & (F \text{ OR } T) \text{ AND } (T \text{ OR } F) \\
 = & T
 \end{aligned}$$

Conclusion:

$$\begin{aligned} & [\text{exists } z: M . p(z, z)] \\ &= P(m1, m1) \text{ OR } P(m2, m2) \\ &= F \text{ OR } F \\ &= F \end{aligned}$$

Express the following sentences in predicate logic. Use types in your formalization. Is the set of formulas consistent? Demonstrate that your answer is correct using the semantics of predicate logic.

All programmer like some computers.

Some programmers use MAC.

Therefore, some people who like some computers use MAC.



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Formalization:

programmer(x) means x is a programmer

usesmac(x) means x uses MAC

likes(x, y) means x likes y

forall x: Person . programmer(x) =>

exists y: Computer . likes(x, y),

exists x: Person . programmer(x) & usesmac(x)

|-

exists x: Person .

(exists y: Computer . likes(x, y) & usesmac(x))

These sentences are *consistent*. Here is an interpretation in which all the formulas are T:

Domain:

People = {John}

Computer = {MacPro}

Mapping:

| Syntax |  | Meaning |
|--------|--|---------|
|--------|--|---------|

|               |  |                      |
|---------------|--|----------------------|
| programmer(.) |  | programmer(John) = T |
|---------------|--|----------------------|

|            |  |                         |
|------------|--|-------------------------|
| likes(.,.) |  | likes(John, MacPro) = T |
|------------|--|-------------------------|

|            |  |                   |
|------------|--|-------------------|
| usesmac(.) |  | usesmac(John) = T |
|------------|--|-------------------|

formula 1:

```
[forall x: Person . programmer(x) =>
  exists y: Computer . likes(x, y)]
= [programmer(^John) =>
  exists y: Computer . likes(^John, y)]
= programmer(John) IMP likes(John, MacPro)
= T IMP T
= T
```

formula 2:

```
[exists x: Person . programmer(x) & usesmac(x)]
= programmer(John) AND usesmac(John)
= T AND T
= T
```

formula 3:

```
[exists x: Person . (exists y: Computer .
  likes(x, y) & usesmac(x))]
= [exists y: Computer .
  likes(^John, y) & usesmac(^John)]
= likes(John, MacPro) AND usesmac(John)
= T AND T
= T
```

If the following arguments are valid, use natural deduction AND semantic tableaux to prove them; otherwise, provide a counterexample.

```
forall x . s(x) | t(x),  
forall x . s(x) => t(x) & k(c, x),  
forall x . t(x) => m(x)  
|-  
m(c)  
where c is a constant
```

## h05q01a (cont'd)

```
#check ND
forall x . s(x) | t(x),
forall x . s(x) => t(x) & k(c, x),
forall x . t(x) => m(x)
|-
m(c)
```

## h05q01a (cont'd)

- 1) forall  $x$  .  $s(x) \mid t(x)$  premise
- 2) forall  $x$  .  $s(x) \Rightarrow t(x) \ \& \ k(c, x)$  premise
- 3) forall  $x$  .  $t(x) \Rightarrow m(x)$  premise
- 4)  $s(c) \mid t(c)$  by forall\_e on 1
- 5)  $s(c) \Rightarrow t(c) \ \& \ k(c, c)$  by forall\_e on 2
- 6)  $t(c) \Rightarrow m(c)$  by forall\_e on 3
- 7) case  $s(c)$  {
  - 8)  $t(c) \ \& \ k(c, c)$  by imp\_e on 5, 7
  - 9)  $t(c)$  by and\_e on 8
  - 10)  $m(c)$  by imp\_e on 6, 9}
- 11) case  $t(c)$  {
  - 12)  $m(c)$  by imp\_e on 6, 11}
- 13)  $m(c)$  by cases on 4, 7-10, 11-12

Is this formula a tautology?

$\vdash (\text{exists } x . p(x)) \Rightarrow \text{forall } y . p(y)$



## h05q01b (cont'd)

No, this formula is not a tautology. Interpretation:

1) Domain = {a, b}

2) Mapping:

| Syntax | Meaning  |
|--------|----------|
| p(.)   | P(a) = T |
|        | P(b) = F |

Conclusion:

$$\begin{aligned} & [(\text{exists } x. p(x)) \Rightarrow \text{forall } y. p(y)] \\ = & (P(a) \text{ OR } P(b)) \text{ IMP } (P(a) \text{ AND } P(b)) \\ = & (T \text{ OR } F) \text{ IMP } (T \text{ AND } F) \\ = & T \text{ IMP } F \\ = & F \end{aligned}$$

Is this argument valid?

```
forall x . p(x) | q(x),  
forall x . !p(x)  
|-  
forall x . q(x)
```

## h05q01d (cont'd)

#check ND

forall x . p(x) | q(x), forall x . !p(x) |- forall x . q(x)

- 1) forall x . p(x) | q(x) premise
- 2) forall x . !p(x) premise
- 3) for every xg {
  - 4) p(xg) | q(xg) by forall\_e on 1
  - 5) case p(xg) {
    - 6) !p(xg) by forall\_e on 2
    - 7) q(xg) by not\_e on 5, 6
  - }
  - 8) case q(xg) {}
  - 9) q(xg) by cases on 4, 5-7, 8-8
- }
- 10) forall x. q(x) by forall\_i on 3-9

# Announcements

- no tutorial next week (Oct 16) (reading week)
- no tutorial the week after (Oct 23) (midterm marking)