

# mmWave Wireless Backhaul Scheduling of Stochastic Packet Arrivals

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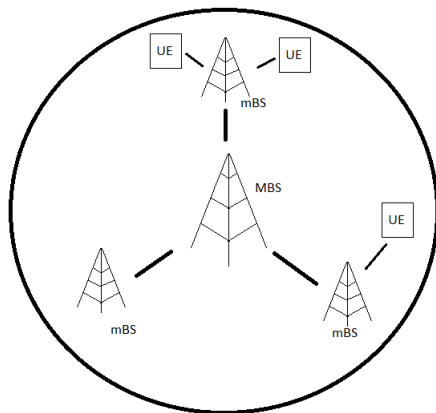
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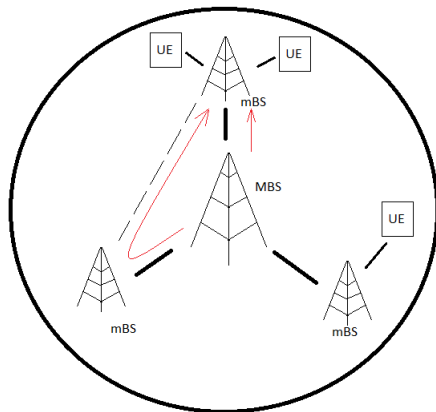
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- 2 Auxiliary definitions
- 3 Results
- 4 Conclusion and Future Work

# Motivation



# Motivation



# Network model

- 1 Macro Base Station (MBS).
- $n$  Micro Base Stations (mBSs).
- Links from MBS to each mBS and from each mBS to every other mBS.
- Time slots, global clock.

# Packet arrivals

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- Expected  $\lambda_i$  packets per round with destination  $i$  (e.g. according to Poisson process).



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- Link  $(i, j)$  can transfer  $c_{i,j}$  packets per round.
- Up to  $k$  links from MBS active.
- Up to 1 link incident with a mBS active.
- Data may be relayed at most once.

# Stochastic stability

## Definition

A mmWBS algorithm is *stochastically stable* against arrival rates  $\lambda$ , if the expected number of rounds to return to a constant-sized queue is finite.

## Related works

- 5G PPP Architecture Working Group, “View on 5G Architecture, version 1.0,” Jul. 2016
- T. Rappaport, S. Sun, R. Mayzus, H. Zhao, Y. Azar, K. Wang, G. Wong, J. Schulz, M. Samimi, and F. Gutierrez, “Millimeter wave mobile communications for 5G cellular: It will work!” IEEE Access, May 2013
- L. Tassiulas and A. Ephremides, “Stability properties of constrained queueing systems and scheduling policies for maximum throughput in multihop radio networks,” Transactions on Automatic Control, 1992

## Our contributions

- Consider mBS as relays to other mBS.
- Show that for “bad” arrival rates no algorithm can be stochastically stable.
- Prove stochastic stability of MaxWeight algorithm for “good” arrival rates.
- Develop a distributed implementation of MaxWeight algorithm that is stochastically stable for “good” arrival rates.

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# Configurations

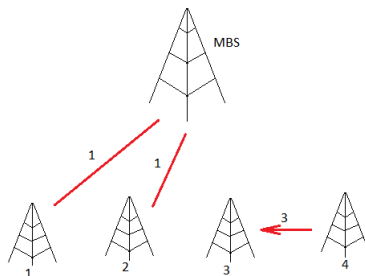
## Definition

A configuration  $N$  is a vector that for every link  $(u, v)$  denotes which packets (i.e., which destination) are to be transmitted (if any).

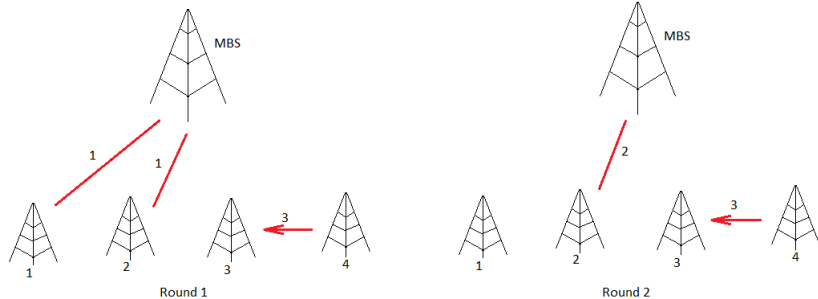
## Definition

A configuration  $N$  is feasible if each mBS has only one incident edge active and there are at most  $k$  edges from MBS active.

# Configurations



# Flow vector



## Stable flow vector

### Definition

A flow vector  $f$  is stable, if:

- it is a linear combination of feasible configurations multiplied by link capacities and
- for each destination  $d$  flow outgoing from MBS equals to flow incoming to mBS  $d$ , which equals arrival rate  $\lambda_d$ .

Intuition: A stable flow vector  $f$  denotes how many packets are transferred over each link per round in a stable execution.

# Bad arrival rates

## Definition

An arrival-rate vector  $\lambda$  is *bad*, if for every combination of feasible configurations  $x$ , and every stable flow vector  $f$ , the following holds:

$$c_{i,j}^{-1} \cdot \sum_d f_{i,j,d} > x_{i,j} \quad \text{for some } x_{i,j} > 0 .$$

## Good arrival rates

### Definition

An arrival-rate vector  $\lambda$  is *good*, if there **exists** combination of feasible configurations  $x$ , and **exists** stable flow vector  $f$ , such that the following holds:

$$c_{i,j}^{-1} \cdot \sum_d f_{i,j,d} < x_{i,j} \quad \text{for all } x_{i,j} > 0 .$$

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# MaxWeight algorithm

At each time step MaxWeight chooses a feasible configuration  $N'$  such that  $N' \leftarrow \arg \max_{N \in \mathcal{S}} \sum_{i,j,d} D_{i,j,d} N_{i,j,d}$ , where

- $D_{i,j,d} = [Q_{i,d} - Q_{j,d}] c_{i,j}$  if  $j \neq i$  and  $D_{i,j,d} = Q_{i,d} c_{i,j}$  otherwise
- $\mathcal{S}$  – set of feasible configurations



## Analysis idea

- Build a Markov chain ( $Q(t)$ ), where  $Q(t)$  denotes a vector of queue sizes  $Q_{i,d}(t)$  at each node  $i$  of each destination  $d$ .

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- Use flows to show that for good arrival rates the drift is negative for almost all starting states  $q$ .
- Use Foster-Lyapunov Theorem to show that some states with small queues are positive recurrent.

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- The chosen configuration is used for the next  $L$  rounds.



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# Conclusion and Future Work

- We have developed an (almost) optimal distributed algorithm for mmWave Wireless Backhaul Scheduling that utilizes relays.
- Open future directions:
  - dynamic link capacities,
  - adversarial packet arrivals,
  - cost of activating/deactivating links.